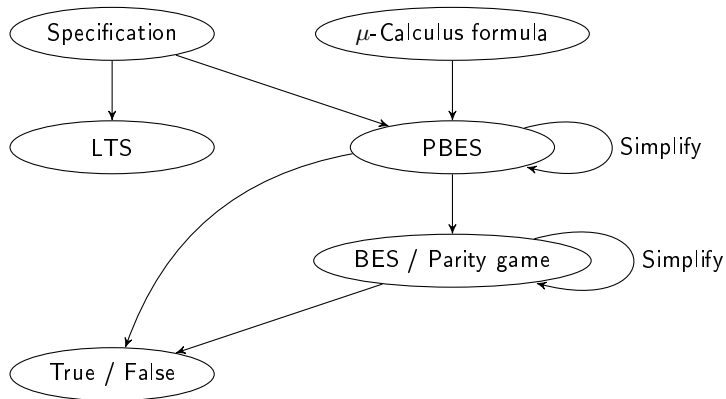


Model checking using Parameterised Boolean Equation Systems

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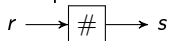


One place buffer



```
sort   Message;  
map   m0 : Message;  
act   r, s : Message;  
proc  P(read : Bool, m : Message) =  
        Σm' : Message.read → r(m').P(¬read, m')  
        ¬read → s(m).P(¬read, m0);  
init  P(true, m0);
```

One place buffer



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sort   Message;
map    m0 : Message;
act    r, s : Message;
proc   P(read : Bool, m : Message) =
          Σ m' : Message. read → r(m'). P(¬read, m')
          ¬read → s(m). P(¬read, m0);
init   P(true, m0);
    
```

“The buffer does not contain any deadlocks”

$$\nu X. [true]X \wedge \langle true \rangle true$$

“The buffer does not contain any deadlocks”

$$\nu X.[true]X \wedge \langle true \rangle true$$

“The buffer does not contain any deadlocks”

$$\nu X.[true]X \wedge \langle true \rangle true$$

```
sort   Message;  
map   m0 : Message;  
pbes  νX(read : Bool, m : Message) =  
        (∀m' : Message.read ⇒ X(¬read, m'))  
        ∧ (¬read ⇒ X(¬read, m0))  
init  X(true, m0)
```

- ▶ Instantiate all data to obtain BES/parity game
 - Reduce modulo equivalences
 - Solve BES/parity game ($NP \cap co - NP$)
- ▶ Simplify PBES on a symbolic level

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Current PBES related research:

- ▶ Allow equivalence reductions on PBES level
- ▶ Improve static simplification techniques
- ▶ Simplify correctness proofs of simplification techniques

Thank you!

6/6

Want to know more? Come talk to me!